ComputerScience





Managing NFV using SDN and Control Theory

What is RINA? [1][2][3]

- RINA: Recursive InterNetwork Architecture
- A clean-slate network architecture that overcomes inherent weaknesses of the current Internet, e.g. security, mobility support
- Based on the fundamental principle that networking is Inter Process Communication (IPC) and only IPC
- Distributed IPC Facility (DIF): a collection of distributed IPC processes with shared states. They provide communication service to application processes over a certain scope (i.e., range of operation)
- Distributed Application Facility (DAF): a set of application processes cooperating to perform a certain function. The function can be a communication service, weather forecast genomics, etc.
- Two design principles: (i) divide and conquer (recursion), and (ii) separation of mechanisms and policies



Fig. 1: RINA overview

Control Theory, SDN and NFV

- Control theory and SDN are key components for NFV deployment
- RINA management architecture is used to manage Virtual Network Function (VNF) instances over the GENI testbed
- Control theory (PI controller) is used to balance load across different VNF instances

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Experiment over GENI

r- °C	 Experimental setup (Fig. 2): two VNF instances running Snort IDS (<i>VNF1</i> and <i>VNF</i>) one OVS switch and one open-flow controller two sources (<i>S1</i> and <i>S2</i>) and one destination (<i>destinat</i>) Traffic is sent to Snort-IDS running on <i>VNF1</i> or <i>VNF2</i> RINA management architecture is used to send load of <i>V</i> instances to controller Load balancer determines the fraction of traffic to divert for VNF1 to VNF2 OVS controller updates openFlow rules on the OVS switch
	Load Balancer
st,	• PI controller (Fig. 3): $x(t) = \max[0, \min[1, x(t-1) + K(\frac{L(t)}{T} - 1)]]$
	<i>x</i> (<i>t</i>): ratio of traffic diverted to VNF2 at time <i>t</i> <i>L</i> (<i>t</i>): load on VNF1 <i>T</i> : target load on VNF1
7	T K X System
	Fig. 3: System load $L(t)$ and target load T of VNF1 is used compute $x(t)$, i.e. ratio of traffic diverted to VNF2
100 80	SNORT-IDS-1 SNORT-IDS-2 80
60 40	MMM Bes 60 MMMM 20 10 10 10 10 10 10 10 10 10 10 10 10 10



Fig. 4: (a) Simple Round Robin load balancing; (b) Load balancing based on PI control (T = 50%)



Management Architecture -2) S2 **S1** tion) VNF1 /NF controller OVS from VNF2 ch destination VNF1 **SNORT IDS RINA Network RINA** App RIB controller DAF **RINA App** Load OVS Balancer controller RIB **CDAP** IDS load info load balancing info → RIB to **RINA** App VNF2 **SNORT IDS** NORT-IDS-1 IORT-IDS-2 Fig. 2: RINA management architecture used for communication between VNF instances and Controller References [1] John Day, Ibrahim Matta and Karim Mattar. "Networking is IPC: A Guiding Principle to a Better Internet". In ReArch 80 100 2008. [2] Boston University RINA Lab. http://csr.bu.edu/rina

[3] Yuefeng Wang, Ibrahim Matta, Flavio Esposito and John Day. "Introducing ProtoRINA: A Prototype for Programming Recursive-Networking Policies." In ACM SIGCOMM CCR, July, 2014.





